CS 267

Tricks with Trees

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Outline

- A log n lower bound to compute any function in parallel
- Reduction and broadcast in O(log n) time
- Parallel prefix (scan) in O(log n) time
- Adding two n-bit integers in O(log n) time
- Multiplying n-by-n matrices in O(log n) time
- Inverting n-by-n triangular matrices in $O(\log^2 n)$ time
- Inverting n-by-n dense matrices in $O(\log^2 n)$ time
- Evaluating arbitrary expressions in O(log n) time
- Evaluating recurrences in O(log n) time

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- Sparse-Matrix-Vector-Multiply (SpMV) using Segmented Scan
- Parallel page layout in a browser (Leo Meyerovich, Ras Bodik)

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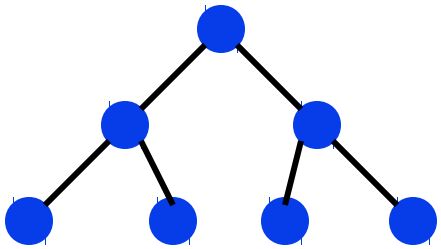
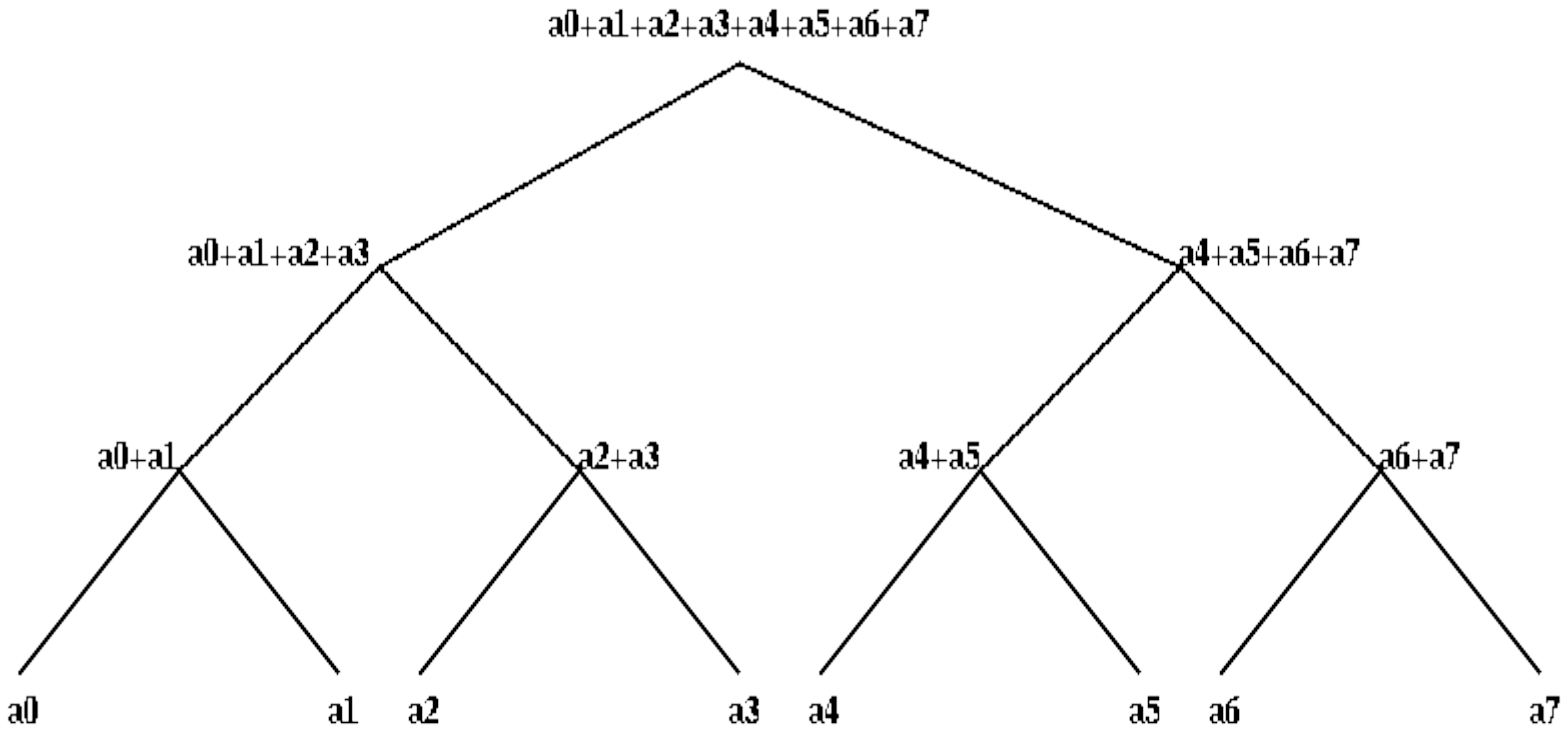
- Sparse-Matrix-Vector-Multiply (SpMV) using Segmented Scan
- Parallel page layout in a browser (Leo Meyerovich, Ras Bodik)
- Solving n-by-n tridiagonal matrices in O(log n) time
- Traversing linked lists
- Computing minimal spanning trees
- Computing convex hulls of point sets

A log n lower bound to compute any function of n variables

- Assume we can only use binary operations, one per time unit
- After 1 time unit, an output can only depend on two inputs
- Use induction to show that after k time units, an output can only depend on 2^k inputs
  - After $\log_2 n$ time units, output depends on at most n inputs
- A binary tree performs such a computation



Broadcasts and Reductions on Trees

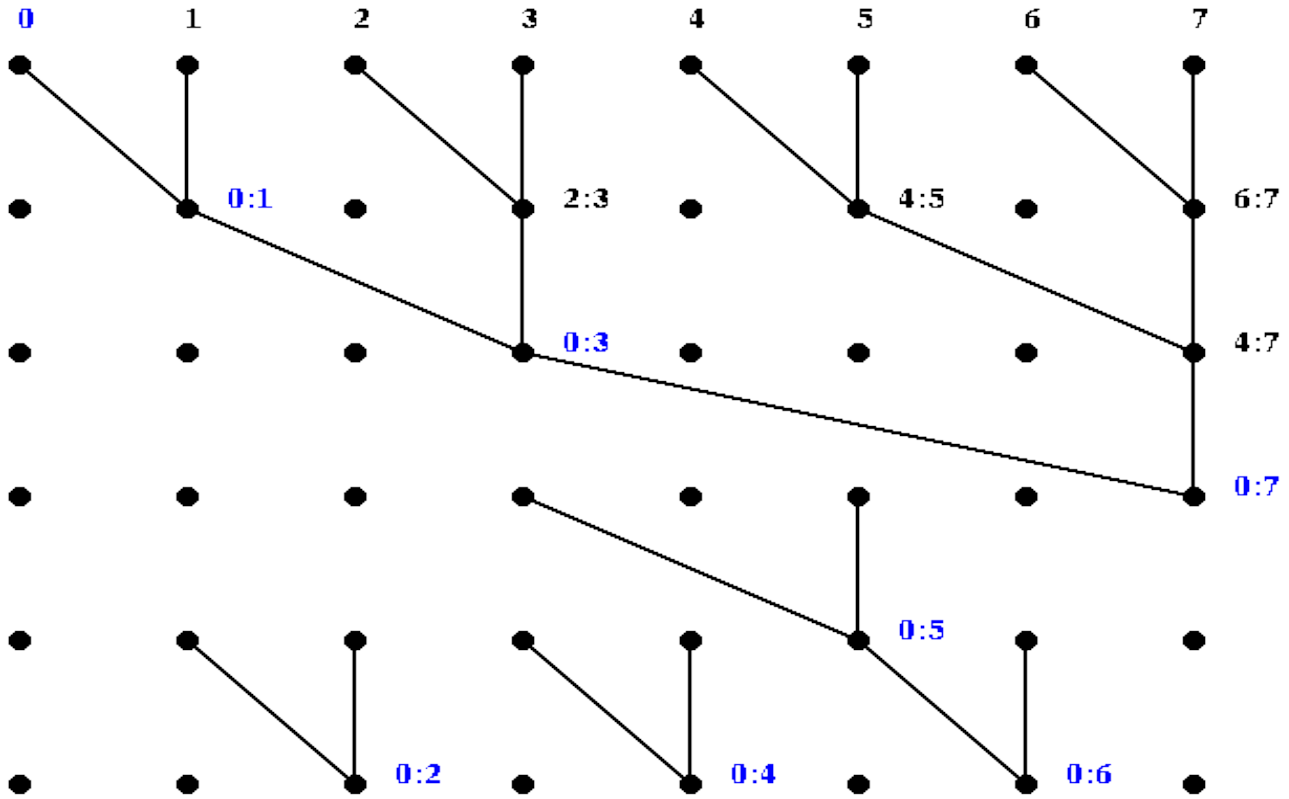


Parallel Prefix, or Scan

- If “+” is an associative operator, and x[0],…,x[p-1] are input data then parallel prefix operation computes

  y[j] = x[0] + x[1] + … + x[j] for j=0,1,…,p-1

- Notation: j:k means x[j]+x[j+1]+…+x[k], blue is final value

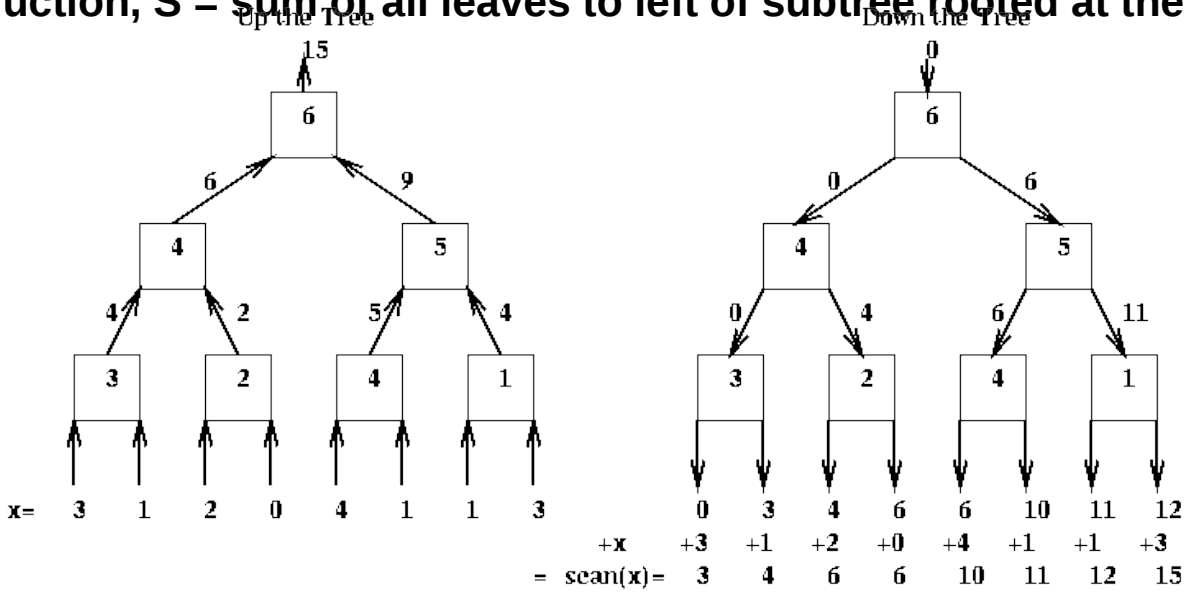


Mapping Parallel Prefix onto a Tree - Details

- Up-the-tree phase (from leaves to root)

1) Get values L and R from left and right children
2) Save L in a local register Lsave
3) Pass sum L+R to parent

- By induction, Lsave = sum of all leaves in left subtree
- Down the tree phase (from root to leaves)
    1) Get value S from parent (the root gets 0)
    2) Send S to the left child
    3) Send S + Lsave to the right child
- By induction, S = sum of all leaves to left of subtree rooted at the parent



E.g., Fibonacci via Matrix Multiply Prefix

$$F_{n+1} = F_n + F_{n-1}$$

$$\begin{pmatrix} F_{n+1} \\ F_n \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F_n \\ F_{n-1} \end{pmatrix}$$

Can compute all F_n by matmul_prefix on

$$\left[\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \right]$$

then select the upper left entry

Adding two n-bit integers in O(log n) time

- **Let a = a[n-1]a[n-2]…a[0] and b = b[n-1]b[n-2]…b[0] be two n-bit binary numbers**
- **We want their sum s = a+b = s[n]s[n-1]…s[0]**

```
c[-1] = 0          … rightmost carry bit
for i = 0 to n-1
    c[i] = ( (a[i] xor b[i])  and  c[i-1] )  or  ( a[i]  and  b[i] )   ... next carry bit
    s[i] = ( a[i] xor b[i] ) xor c[i-1]
```

- **Challenge: compute all c[i] in O(log n) time via parallel prefix**

for all (0 <= i <= n-1) p[i] = a[i] xor b[i] … propagate bit
for all (0 <= i <= n-1) g[i] = a[i] and b[i] … generate bit

$$\begin{pmatrix} c[i] \\ 1 \end{pmatrix} = \begin{pmatrix} (\, p[i] \text{ and } c[i-1] \,) \text{ or } g[i] \\ 1 \end{pmatrix} = \begin{pmatrix} p[i] & g[i] \\ 0 & 1 \end{pmatrix} * \begin{pmatrix} c[i-1] \\ 1 \end{pmatrix} = C[i] * \begin{pmatrix} c[i-1] \\ 1 \end{pmatrix}$$

… 2-by-2 Boolean matrix multiplication (associative)

$$= C[i] * C[i-1] * \ldots C[0] * \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

… evaluate each P[i] = C[i] * C[i-1] * … * C[0] by parallel prefix

- **Used in all computers to implement addition - Carry look-ahead**

Other applications of scans

- There are many applications of scans, some more obvious than others
  - add multi-precision numbers (represented as array of numbers)
  - evaluate recurrences, expressions
  - solve tridiagonal systems (numerically unstable!)
  - implement bucket sort and radix sort
  - to dynamically allocate processors
  - to search for regular expression (e.g., grep)
- Names: +\ (APL), cumsum (Matlab), MPI_SCAN
- Note: 2n operations used when only n-1 needed

Multiplying n-by-n matrices in O(log n) time

- For all (1 <= i,j,k <= n) P(i,j,k) = A(i,k) * B(k,j)
  - cost = 1 time unit, using n^3 processors
- For all (1 <= i,j <= n) $C(i,j) = \sum_{k=1}^{n} P(i,j,k)$
  - cost = O(log n) time, using a tree with $n^3 / 2$ processors

Inverting triangular n-by-n matrices in $O(\log^2 n)$ time

° **Fact:**

$$\begin{pmatrix} A & 0 \\ C & B \end{pmatrix}^{-1} = \begin{pmatrix} A^{-1} & 0 \\ -B^{-1}CA^{-1} & B^{-1} \end{pmatrix}$$

° **Function Tri_Inv(T)** … assume $n = \dim(T) = 2^m$ for simplicity

```
If T is 1-by-1
   return 1/T
else
   … Write T = [A 0; C B]
   In parallel do {
        invA = Tri_Inv(A)
        invB = Tri_Inv(B)  }      … implicitly uses a tree
   newC = -invB * C * invA
   Return [invA 0; newC invB]
```

° **time(Tri_Inv(n)) = time(Tri_Inv(n/2)) + O(log(n))**

- **Change variable to m = log n to get time(Tri_Inv(n)) = $O(\log^2 n)$**

Inverting Dense n-by-n matrices in $O(\log^2 n)$ time

- **Lemma 1: Cayley-Hamilton Theorem**
  - expression for A^{-1} via characteristic polynomial in A
- **Lemma 2: Newton's Identities**
  - Triangular system of equations for coefficients of characteristic polynomial, where matrix entries = s_k
- **Lemma 3:** $s_k = \text{trace}(A^k) = \sum_{i=1}^{n} A^k[i,i] = \sum_{i=1}^{n} [\lambda_i(A)]^k$
- **Csanky's Algorithm (1976)**
  1) Compute the powers $A^2, A^3, \ldots, A^{n-1}$ by parallel prefix
     cost = $O(\log^2 n)$
  2) Compute the traces $s_k = \text{trace}(A^k)$
     cost = $O(\log n)$
  3) Solve Newton identities for coefficients of characteristic polynomial
     cost = $O(\log^2 n)$
  4) Evaluate A^{-1} using Cayley-Hamilton Theorem
     cost = $O(\log n)$
  - **Completely numerically unstable**

Evaluating arbitrary expressions

- Let E be an arbitrary expression formed from +, -, *, /, parentheses, and n variables, where each appearance of each variable is counted separately
- Can think of E as arbitrary expression tree with n leaves (the variables) and internal nodes labeled by +, -, * and /
- Theorem (Brent): E can be evaluated in O(log n) time, if we reorganize it using laws of commutativity, associativity and distributivity
- Sketch of (modern) proof: evaluate expression tree E greedily by repeatedly
  - collapsing all leaves into their parents at each time step
  - evaluating all "chains" in E with parallel prefix

Evaluating recurrences

- **Let $x_i = f_i(x_{i-1})$, f_i a rational function, x_0 given**
- **How fast can we compute x_n?**
- **Theorem (Kung): Suppose degree(f_i) = d for all i**
  - **If d=1, x_n can be evaluated in O(log n) using parallel prefix**
  - **If d>1, evaluating x_n takes $\Omega(n)$ time, i.e. no speedup is possible**
- **Sketch of proof when d=1**

  $x_i = f_i(x_{i-1}) = (a_i * x_{i-1} + b_i)/(c_i * x_{i-1} + d_i)$ can be written as

  $$x_i = \frac{num_i}{dem_i} = \frac{(a_i * num_{i-1} + b_i * den_{i-1})}{(c_i * num_{i-1} + d_i * den_{i-1})} \text{ or}$$

  $$\begin{bmatrix} num_i \\ dem_i \end{bmatrix} = \begin{bmatrix} a_i & b_i \\ c_i & d_i \end{bmatrix} * \begin{bmatrix} num_{i-1} \\ den_{i-1} \end{bmatrix} = M_i * \begin{bmatrix} num_{i-1} \\ den_{i-1} \end{bmatrix} = M_i * M_{i-1} * \ldots * M_1 * \begin{bmatrix} num_0 \\ den_0 \end{bmatrix}$$

  Can use parallel prefix with 2-by-2 matrix multiplication
- **Sketch of proof when d>1**
  - **degree(x_i) as a function of x_0 is d^i**
  - **After k parallel steps, degree(anything) ≤ 2^k**
  - **Computing x_i take $\Omega(i)$ steps**

Image Segmentation (1/4)

- Contours are subjective – they depend on perspective
  - Surprise: Humans agree (somewhat)
- Goal: generate contours automatically
  - Use them to break images into separate segments (subimages)
  - J. Malik's group has leading algorithm
  - Enable automatic image search and retrieval ("Find all the pictures with Fred")

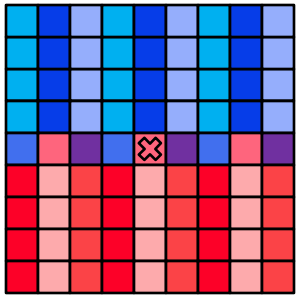
Image

Human Generated Contours

Machine Generated Contours

Image Segmentation (2/4)

- Think of image as matrix A(i,j) of pixels
  - Each pixel has separate R(ed), G(reen), B(lue) intensities
- Bottleneck (so far) of Malik's algorithm is to compute other matrices indicating whether pixel (i,j) likely to be on contour
  - Ex: C(i,j) = average "R intensity" of pixels in rectangle above (i,j) – average "R intensity" of pixels in rectangle below (i,j)
  - C(i,j) large for pixel (i,j) marked with ⊠, so (i,j) likely to be on contour

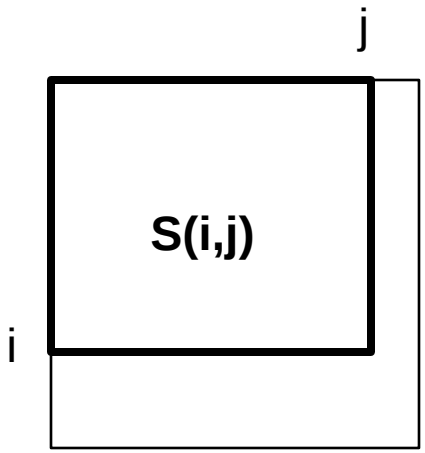


- Algorithm eventually computes eigenvectors of sparse matrix with entries computed from matrices like C
  - Analogous to graph partitioning in later lecture

Image Segmentation (3/4)

- **Bottleneck: Given A(i,j), compute C(i,j) where**
  - Sa(i,j) = sum of A(i,j) for entries in k x (2k+1) rectangle above A(i,j)

    $= \Sigma$ A(r,s) for $i-k \le r \le i-1$ and $j-k \le s \le j+k$
  - Sb(i,j) = similar sum of rectangle below A(i,j)
  - C(i,j) = Sa(i,j) – Sb(i,j)
- **Approach (Bryan Catanzaro)**
  - Compute S(i,j) = Σ A(r,s) for $r \le i$ and $s \le j$
  - Then sum of A(i,j) over any rectangle ($I_{low} \le i \le I_{high}$, $J_{low} \le j \le J_{high}$) is $S(I_{high}, J_{high}) - S(I_{low}-1, J_{high}) - S(I_{high}, J_{low}-1) + S(I_{low}-1, J_{low}-1)$



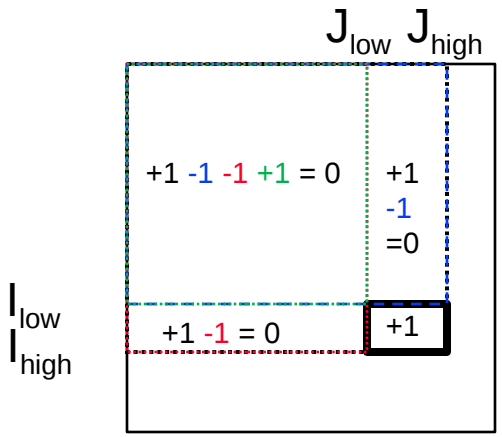
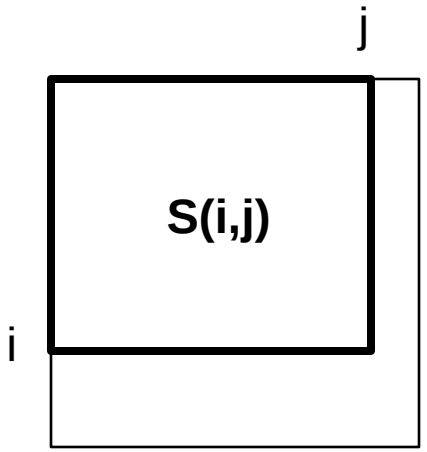


Image Segmentation (4/4)

- New Bottleneck: Given A(i,j), compute S(i,j) where
  - S(i,j) = Σ A(r,s) for r ≤ i and s ≤ j
- "2 dimensional parallel prefix"
  - Do parallel prefix independently on each row of A(i,j) :
    - $S_{row}(i,j) = \Sigma A(i,s)$ for s ≤ j
  - Do parallel prefix independently on each column of S_{row}
    - $S(i,j) = \Sigma S_{row}(r,j)$ for r ≤ i $= \Sigma A(r,s)$ for s ≤ j and r ≤ i



Sparse-Matrix-Vector-Multiply (SpMV) y = A*x Using Segmented Scan (SegScan)

- **Segscan computes prefix sums of arbitrary segments**

  Segscan ([3, 1, 4, 5, 6, 1, 2, 3],
  [T, F, F, T, T, F, F, T])
  = [3, 4, 8, 5, 6, 7, 9, 3]

- **Use CSR format of Sparse Matrix A, store x densely**

$$A = \begin{bmatrix} 0 & 2 & 3 & 0 \\ 4 & 0 & 0 & 5 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Val = [1 2 3 2 4 5 3 1]
Col_Ind = [1 3 4 1 2 5 1 5]
Row_Ptr=[1 4 7 9]

$$x = \begin{bmatrix} 7 \\ 8 \\ 2 \\ 1 \\ 3 \end{bmatrix}$$

- **Create array P of all nonzero A(i,j)*x(j) = Val(k)*x(Col_Ind(k))**

  P = [7 4 3 14 32 15 21 3]

- **Create array S showing where segments (rows) start**

  S = [1 0 0 1 0 0 1 0]

- **Compute SegScan(P, S) = [7 11 14 14 46 61 21 24]**

- **Extract A*x = [14 61 24]**

- **www.cs.cmu.edu/afs/cs.cmu.edu/project/scandal/public/papers/CMU-CS- 93-173.ps.Z**

Page layout in a browser

- Applying layout rules to html description of a webpage is a bottleneck, scan can help
- Simplest example
  - Given widths [x_1, x_2, … , x_n] of items to display on page, where should each item go?
  - Item j starts at $x_1 + x_2 + \ldots + x_{j-1}$
- Real examples have complicated constraints
  - Defined by general trees, since in html each object to display can be composed of other objects
  - To get location of each object, need to do preorder traversal of tree, "adding up" constraints of previous objects
  - Scan can do preorder traversal of any tree in parallel
    - Not just binary trees
- Ras Bodik, Leo Meyerovich

Summary of tree algorithms

- Lots of problems can be done quickly - in theory - using trees
- Some algorithms are widely used
  - broadcasts, reductions, parallel prefix
  - carry look ahead addition
- Some are of theoretical interest only
  - Csanky's method for matrix inversion
  - Solving general tridiagonals (without pivoting)
  - Both numerically unstable
  - Csanky needs too many processors
- Embedded in various systems
  - MPI, Split-C, Titanium, NESL, other languages
  - CM-5 hardware control network